

APPENDIX B

MC68EC040

NOTE

Rev. 2.2 contains timing information for 40 MHz operation. Refer to chang bars. Some TBD values will be filled in shortly.

All references to MC68EC040 also apply to the MC68EC040V. Refer to **Appendix C MC68040V and MC68EC040V** for more information on the MC68EC040V.

The MC68EC040 is Motorola's third generation of M68000-compatible, high-performance, 32-bit microprocessors. The MC68EC040 is an embedded controller employing a highly integrated architecture to provide very high performance in a monolithic HCMOS device. The MC68EC040 integrates an MC68040-compatible integer unit, an access control unit (ACU), and independent 4-Kbyte instruction and data caches. A six-stage instruction pipeline, multiple internal buses, and a full internal Harvard architecture, including separate caches for both instruction and data accesses, provides a high degree of instruction execution parallelism. The inclusion of on-chip bus snooping logic, which directly supports cache coherency in multimaster applications, enhances cache functionality.

The MC68EC040 is user-object-code compatible with previous members of the M68000 family and is specifically optimized to reduce the execution time of compiler-generated code. The MC68EC040 is pin compatible with the MC68040 and MC68LC040. The MC68EC040 is implemented in Motorola's latest HCMOS technology, providing an ideal balance between speed, power, and physical device size. Figure B-1 provides a simplified block diagram of the MC68EC040.

The main features of the MC68EC040 include:

- MC68040-Compatible Integer Execution Unit
- 4-Kbyte Instruction Cache and 4-Kbyte Data Cache Accessible Simultaneously
- 32-Bit, Nonmultiplexed External Address and Data Buses with Synchronous Bursting Interface
- User-Object-Code Compatible with All M68000 Microprocessors
- Concurrent Integer Unit, ACU, and Bus Controller Operation Maximizes Throughput
- Low-Latency Bus Accesses for Reduced Cache-Miss Penalty
- Multimaster/Multiprocessor Support via Bus Snooping
- 4-Gbyte Direct Addressing Range

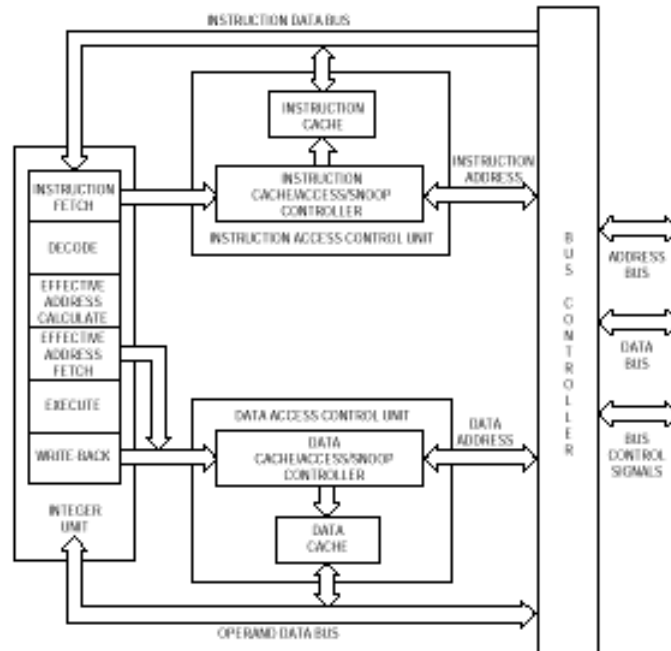


Figure B-1. MC68EC040 Block Diagram

With the exception of the memory management unit (MMU), the floating-point unit (FPU), and their respective registers, the MC68EC040 programming model, data formats and types, instruction set (except all instructions beginning with an “F”, PTEST, and PFLUSH), and caches are the same as described in **Section 1 Introduction** for the MC68040. Figures B-2 and B-3 illustrate the programming model and functional signal groups for the MC68EC040.

B.1 MC68EC040 DIFFERENCES

The following differences exist between the MC68EC040 and MC68040:

- Two independent access control units (ACUs) replace the MC68040 MMUs. The ACU has four corresponding registers (access control registers) that the MC68040 implements as data transparent translation registers. The page size is fixed at 4 Kbytes.
- PTEST and PFLUSH instructions cause an indeterminate result (i.e., an undetermined number of bus cycles); the user should not execute them on the MC68EC040.
- The MC68EC040 does not contain an FPU which causes unimplemented floating-point exceptions to occur using a new stack frame format.

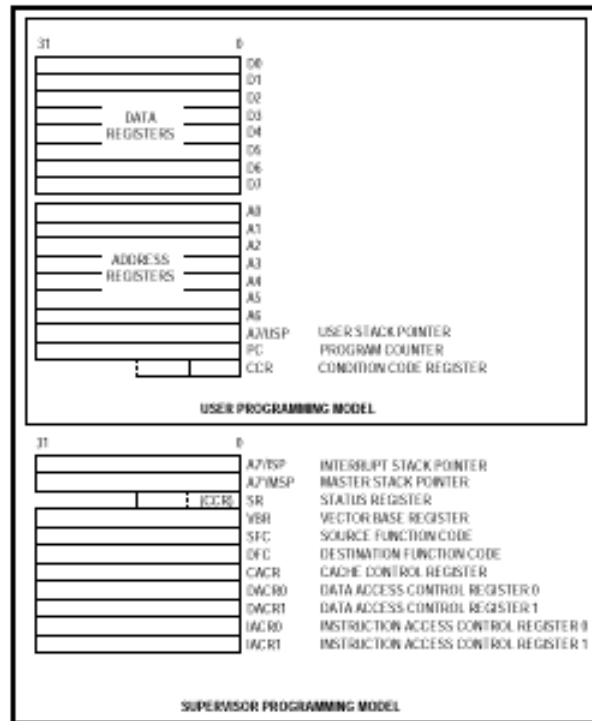


Figure B-2. MC68EC040 Programming Model

- The DLE and $\overline{\text{MDIS}}$ pin names have been changed to JS0 and JS1, respectively.
- The MC68EC040 does not implement the DLE mode, multiplexed, or output buffer impedance selection modes of operation. The MC68EC040 implements only the small output buffer mode of operation. All timing and drive capabilities of the MC68EC040 are equivalent to those of the MC68040 in the small buffer mode of operation.

B.2 JTAG SCAN (JS1–JS0)

The MC68040 $\overline{\text{MDIS}}$ and DLE pin names have been changed to JS1 and JS0 respectively. During normal operation, the JS1 and JS0 pin cannot float, they must be tied to GND or V_{cc} directly or through a resistor. During board testing, these pins retain the functionality of the JTAG scan of the MC68040 for compatibility purposes. Refer to **Section 6 IEEE 1149.1A Test Access Port (JTAG)** for details concerning IEEE 1149.1 *Standard Test Access Port and Boundary Scan Architecture*.

B.3 ACCESS CONTROL UNITS

The information in this section replaces the information in **Section 3 Memory Management Unit (Except MC68EC040 and MC68EC040V)**. When reading **Section 4 Instruction and Data Caches**, disregard any references to the MMU; remember the functionality of the access control registers has replaced that of transparent translation registers. The

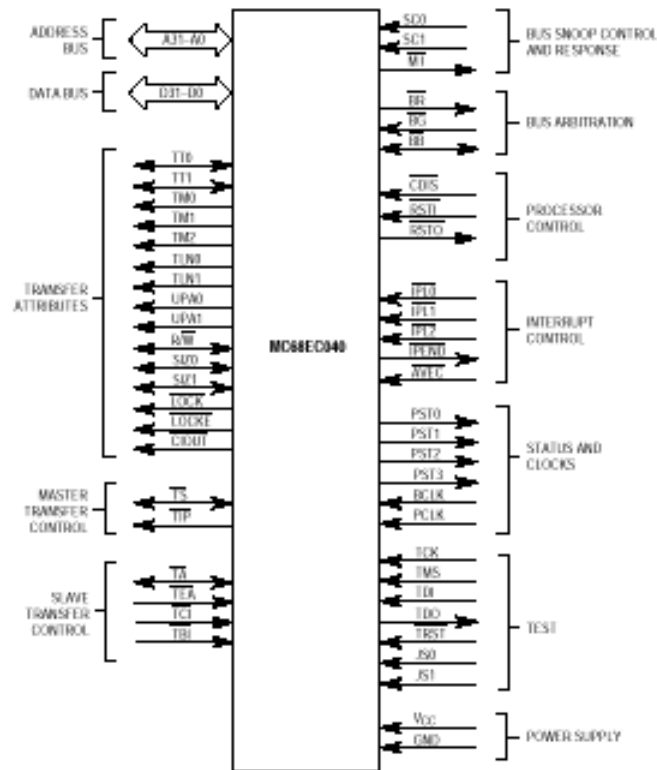


Figure B-3. MC68EC040 Functional Signal Groups

MC68EC040 contains two independent ACUs, one for instructions and one for data. Each ACU allows memory selections to be made requiring attributes particular to peripherals, shared memory, or other special memory requirements. The following paragraphs describe the ACUs and the access control registers contained in them.

B.3.1 Access Control Registers

Each ACU has two independent access control registers (ACRs). The instruction ACU contains the instruction access control registers (IACR0 and IACR1). The data ACU contains the data access control registers (DACR0 and DACR1). Both ACRs provide and control status information for access control of memory in the MC68EC040. Only programs that execute in the supervisor mode using the MOVEC instruction can directly access the ACRs.

The 32-bit ACRs each define blocks of MC68EC040:address space for access control. These blocks of address space can overlap or be separate, and are a minimum of 16 Mbytes. Three blocks are used with two user-defined attributes, cachability control and optional write protection. The ACRs specify a block of address space as serialized noncachable for peripheral selections and as write-through for shared blocks of address space in multi-processing applications. The ACRs can be configured to support many embedded

control applications while maintaining cache integrity. Refer to **Section 4 Instruction and Data Caches** for details concerning cachability. Figure B-4 illustrates the ACR format.

31	24	23	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
LOGICAL ADDRESS BASE		LOGICAL ADDRESS MASK		E	S	0	0	0	U1	U0	0	CM	0	0	W	0	0		

Figure B-4. MC68EC040 Access Control Register Format

ADDRESS BASE

This 8-bit field is compared with physical address bits A31–A24. Addresses that match in this comparison (and are otherwise eligible) are accessible.

ADDRESS MASK

This 8-bit field contains a mask for the ADDRESS BASE field. Setting a bit in the ADDRESS MASK field causes the processor to ignore the corresponding bit in the ADDRESS BASE field. Setting some of the ADDRESS MASK bits to ones obtains blocks of memory larger than 16 Mbytes. The low-order bits of this field are normally set to define contiguous blocks larger than 16 Mbytes, although contiguous blocks are not required.

E—Enable

This bit enables and disables transparent translation of the block defined by this register. Refer to **Section 3 Memory Management Unit (Except MC68EC040 and MC68EC040V)** for details on transparent translation.

- 0 = Access control disabled.
- 1 = Access control enabled.

S—Supervisor/User Mode

This field specifies the way FC2 is used in matching an address:

- 00 = Match only if FC2 = 0 (user mode access).
- 01 = Match only if FC2 = 1 (supervisor mode access).
- 10, 11 = Ignore FC2 when matching.

U1, U0—User Page Attributes

These two bits drive on the user page attribute signals (UPA1 and UPA0). If an external bus transfer results from the access, U0 and U1 are echoed to the UPA0 and UPA1 signals, respectively. The user can program these bits to support extended addressing, bus snooping, or other applications. The MC68EC040 does not interpret these bits.

CM—Cache Mode

This field selects the cache mode and access serialization for a page as follows:

- 00 = Cachable, Write-through
- 01 = Cachable, Copyback
- 10 = Noncachable, Serialized
- 11 = Noncachable

Detailed information on caching modes is available in **Section 4 Instruction and Data Caches**, and information on serialization is available in **Section 7 Bus Operation**.

W—Write Protect

This bit indicates if the transparent block is write protected. If set, write and read-modify-write accesses are aborted as if the R-bit in a table descriptor were clear. Refer to **3.2.2 Descriptors** for a description of table descriptors.

- 0 = Read and write accesses permitted.
- 1 = Write accesses not permitted.

B.3.2 Address Comparison

The following description of address comparison assumes that the ACRs are enabled. Clearing the E-bit in each ACR independently disables access control, causing the processor to ignore it.

When an ACU receives a physical address, the privilege mode and the eight high-order bits of the address are compared to the block of addresses defined by the two ACRs for the corresponding ACU. Each block of address space for an ACR contains an S-field, a BASE ADDRESS field, and an ADDRESS MASK field. The S-field allows for matching either user or supervisor accesses (or both). Setting a bit in the ADDRESS MASK field causes the corresponding bit of the ADDRESS BASE to be ignored in the address comparison and privilege mode. Setting successively higher order bits in the ADDRESS MASK field increases the size of the block of address space.

The address for the current bus cycle and an ACR address match when the privilege mode and address bits for each (not including the masked bits) are equal. Each ACR specifies write protection for the block of address space. Enabling write protection for a block of address space causes the abortion of write or read-modify-write accesses to the block, and an access error exception occurs.

By appropriately configuring an ACR, flexible mappings can be specified. For example, to control access to the user address space, the S-field equals \$0, and the ADDRESS MASK field equals \$FF in all four ACRs. To control access to the supervisor address space (\$00000000–\$0FFFFFFF) with write protection, the BASE ADDRESS field = \$0X, the ADDRESS MASK field equals \$0F, the W-bit is set to one, and the S-field = \$1. The inclusion of independent ACRs in both the instruction ACU (IACU) and data ACU (DACU) provides an exception to the merged instruction and data address space, allowing different access control for instruction and operand accesses. Also, since the instruction memory unit is only used for instruction prefetches, different instruction and data ACRs can cause PC relative operand fetches to be translated differently from instruction prefetches.

Matching either of the ACRs in a corresponding ACU during an access to a memory unit completes the access with the ACU. If both registers match, the access uses the xACR0 status bits. Addresses are passed through without translation if there is no match in the ACRs and no table search occurs. The MC68EC040 does not perform table searches.

B.3.3 Effect of $\overline{\text{RSTI}}$ on the ACU

When the assertion of the reset input ($\overline{\text{RSTI}}$) signal resets the MC68EC040, the E-bits of the ACRs are cleared, disabling address access control.

B.4 SPECIAL MODES OF OPERATION

This part of the M68040 User's Manual does not apply to the MC68EC040. The MC68EC040 does not sample the $\overline{\text{IPL2}}\text{--}\overline{\text{IPL0}}$, $\overline{\text{CDIS}}$, JS0 (DLE on the MC68040), or JS1 ($\overline{\text{MDIS}}$ on the MC68040) pins on the rising edge of $\overline{\text{RSTI}}$.

An external device asserts $\overline{\text{RSTI}}$ to reset the processor. When power is applied to the system, external circuitry should assert $\overline{\text{RSTI}}$ for a minimum of 10 BCLK cycles after V_{CC} is within tolerance. Figure B-5 is a functional timing diagram of the power-on reset operation, illustrating the relationships between V_{CC} , $\overline{\text{RSTI}}$, and bus signals. The BCLK and PCLK clock signals are required to be stable by the time V_{CC} reaches the minimum operating specification. $\overline{\text{RSTI}}$ is internally synchronized for two BCLKS before being used, and must meet the specified setup and hold times to BCLK (specifications #51 and #52 in **MC68EC040 Electrical Characteristics**) only if recognition by a specific BCLK rising edge is required.

Once $\overline{\text{RSTI}}$ is negated, the processor is internally held in reset for another 128 clock cycles. During the reset period, all three-statable signals are three-stated, and the rest are driven to their inactive state. Once the internal reset signal negates, all bus signals remain in a high-impedance state until the processor is granted the bus. After this, the first bus cycle for reset exception processing begins. In Figure B-6, the processor assumes implicit ownership of the bus before the first bus cycle begins. The levels on the $\overline{\text{CDIS}}$, JS1 ($\overline{\text{MDIS}}$ on the MC68040), and $\overline{\text{IPL2}}\text{--}\overline{\text{IPL0}}$ signals are not sampled when $\overline{\text{RSTI}}$ is negated.

For processor resets after the initial power-on reset, should be asserted for at least 10 clock periods. Figure B-6 illustrates timing associated with a reset when the processor is executing bus cycles. Note that $\overline{\text{BB}}$ and $\overline{\text{TIP}}$ (and $\overline{\text{TA}}$ driven during a snooped access) are asserted before transitioning to a three-state level. Processor reset causes any bus cycle in progress to terminate as if $\overline{\text{TA}}$ or $\overline{\text{TEA}}$ had been asserted. Also, the processor initializes registers appropriately for a reset exception.

When a RESET instruction is executed, the processor drives the reset out ($\overline{\text{RSTO}}$) signal for 512 BCLK cycles. In this case, the processor resets the external devices of the system, and the internal registers of the processor are unaffected. The external devices connected to $\overline{\text{RSTO}}$ are reset at the completion of the RESET instruction. An $\overline{\text{RSTI}}$ signal that is asserted to the processor during execution of a RESET instruction immediately resets the processor and causes $\overline{\text{RSTO}}$ to negate. $\overline{\text{RSTO}}$ can be logically ANDed with the external signal driving

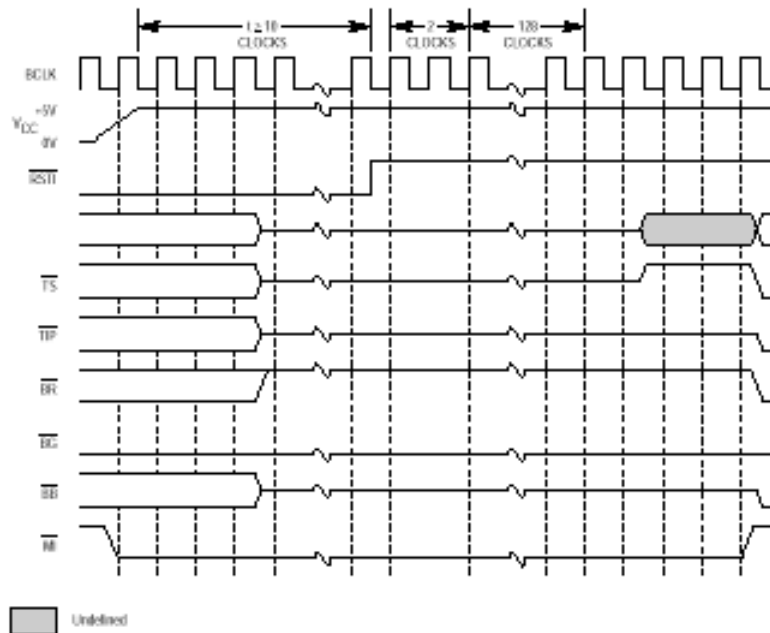


Figure B-5. MC68EC040 Initial Power-On Reset Timing

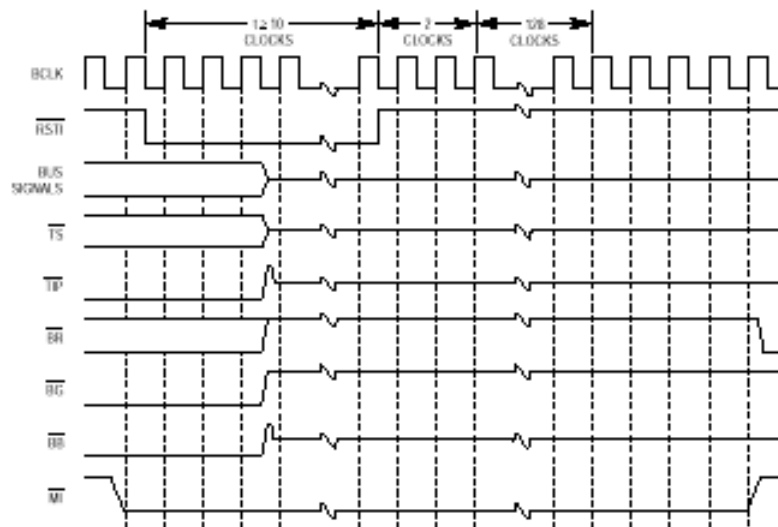


Figure B-6. MC68EC040 Normal Reset Timing

\overline{RSTI} to derive a system reset signal that is asserted for both an external processor reset and execution of a RESET instruction.

B.5 EXCEPTION PROCESSING

The MC68EC040 provides five different stack frames for exception processing and allows for a MC68040-specific stack frame. Refer to **Section 8 Exception Processing** for details on exception processing.

B.5.1 Unimplemented Floating-Point Instructions and Exceptions

All legal MC68040 and MC68881/MC68882 floating-point instructions are defined as unimplemented floating-point instructions on the MC68EC040. These instructions generate an eight-word stack frame (format \$4) during exception processing before taking an F-line exception. These instructions trap as an F-line exception and can be emulated in software by the F-line exception handler to maintain user-object-code compatibility.

The MC68EC040 assists the emulation process by distinguishing unimplemented floating-point instructions from other unimplemented F-line instructions. To aid emulation, the effective address is calculated and saved in the format \$4 stack frame. This simplifies and speeds up the emulation process by eliminating the need for the emulation routine to determine the effective address and by providing information required to emulate the instruction on the exception stack frame in the supervisor address space. However, the floating-point instruction can reside in user space; therefore, the floating-point unimplemented exception handler may need to access user instruction space. The following processing steps occur for an unimplemented floating-point instruction:

1. When an unimplemented floating-point instruction is encountered, the instruction is partially decoded, and the effective address is calculated, if required.
2. The processor waits for all previous integer instructions, write-backs, and associated exception processing to complete before beginning exception processing for the unimplemented floating-point instruction. Any access error that occurs in completing the write-backs causes an access error exception, and the resulting stack frame indicates a pending unimplemented floating-point instruction exception. The access error exception handler then completes the write-backs in software, and exception processing for the unimplemented floating-point instruction exception begins immediately after return from the access error handler.
3. The processor begins exception processing for the unimplemented floating-point instruction by making an internal copy of the current SR. The processor then enters the supervisor mode and clears the trace bits (T1 and T0). It creates a format \$4 stack frame and saves the internal copy of the SR, PC, vector offset, calculated effective address, and PC value of the faulted instruction in the stack frame.

The effective address field of the format \$4 stack frame contains the calculated effective address of the operand for the faulted floating-point instruction using the addressing mode in which the effective address is calculated. For immediate and register direct addressing modes, this field is \$0. The saved PC value is the logical address of the instruction that follows the unimplemented floating-point instruction. This value will be restored during RTE execution. The vector offset format number (\$4) is used for this eight-word stack frame. Note that an MC68040 cannot correctly handle a stack format \$4. The PC of the faulted instruction contains a long-word PC of the floating-point instruction that caused the trap to occur. The information is provided so that the in-

struction is available for software emulation of floating-point instructions. The processor generates exception vector number 11 for the unimplemented F-line instruction exception vector, fetches the address of the F-line exception handler from the exception vector table, and begins execution of the handler after prefetching instructions to fill the pipeline. Refer to **Section 8 Exception Processing** for details about exception processing.

B.5.2 MC68EC040 Stack Frames

When the processor executes an RTE instruction, it examines the stack frame on top of the active supervisor stack to determine if it is a valid frame and what type of context restoration it requires. The set of stack frames included for exception processing are four- and six-word stack frames, a four-word throwaway stack frame, an access error stack frame, and a new eight-word unimplemented floating-point stack frame. The stack frame that the MC68040 can generate and the MC68EC040 can process is the floating-point post-instruction stack frame. Refer to **Section 8 Exception Processing** for details about exception stack frames.

Eight-Word Stack Frame (Format \$4)

Stack Frames	Exception Types	Stacked PC Points To
(UNABLE TO LOCATE ART)	<ul style="list-style-type: none"> The MC68040 cannot generate or read this stack. 	<ul style="list-style-type: none"> Effective address field is the address of the faulted instruction operand.

When the MC68EC040 writes or reads a stack frame, it uses long-word operand transfers wherever possible. Using a long-word-aligned stack pointer greatly enhances exception processing performance. The processor does not necessarily read or write the stack frame data in sequential order. The system software should not depend on a particular exception generating a particular stack frame. For compatibility with future devices, the software should be able to handle any type of stack frame for any type of exception. The MC68EC040 does not generate the floating-point post-instruction stack frame. The MC68040 cannot accept the eight-word unimplemented stack frame. The MC68EC040 can handle all MC68040 stack frame formats.

B.6 SOFTWARE CONSIDERATIONS

The following MC68EC040 instructions are different from the MC68040: PTEST, PFLUSH, CPUSH, CINV, MOVEC, and all floating-point instructions.

The PTEST and PFLUSH instructions should not be executed. Execution of the PTEST instruction causes random bus cycles to occur. Execution of the PFLUSH instruction produces indeterminate results. Neither instruction causes the MC68EC040 to generate an exception.

The CPUSH and CINV instructions require special consideration. A page is defined as a 4-Kbyte block of external memory. The CPUSH and CINV page instruction opcodes can be used to push or invalidate 4-Kbyte blocks of memory. The MC68EC040 does not support 8-Kbyte pages.

The MOVEC to URP and SRP instructions are not valid and will produce indeterminate results. Each ACU has a status register and translation control register that replace the MMU

status register and translation control register of the MC68040. The MMU status register opcode of the MOVEC instruction can modify the ACU status register. The MC68EC040 ACU status register does not provide additional functionality to the ACU and is only provided for compatibility with the ACU MC68EC030 status register. The ACU status register may not be implemented in future M68EC0X0 products.

B.7 MC68EC040 ELECTRICAL CHARACTERISTICS

The following paragraphs provide information on the maximum rating and thermal characteristics for the MC68EC040 only. Refer to Appendix C MC68040V and MC68EC040V for more information on electrical characteristics for the MC68EC040V. This section is subject to change. For the most recent specifications, contact a Motorola sales office or complete the registration card at the end of this manual.

B.7.1 Maximum Ratings

Table 12-2.

Characteristic	Symbol	Value	Unit
Supply Voltage	V_{CC}	-0.3 to +7.0	V
Input Voltage	V_{in}	-0.5 to +7.0	V
Maximum Operating Junction Temperature	T_J	110	°C
Minimum Operating Ambient Temperature	T_A	0	°C
Storage Temperature Range	T_{stg}	-55 to 150	°C

This device contains protective circuitry against damage due to high static voltages or electrical fields; however, it is advised that normal precautions be taken to avoid application of any voltages higher than maximum-rated voltages to this high-impedance circuit. Reliability of operation is enhanced if unused inputs are tied to an appropriate logic voltage level (e.g., either GND or V_{CC}).

B.7.2 Thermal Characteristics

Characteristic	Symbol	Value	Rating
Thermal Resistance, Junction to Case— PGA Package	θ_{JC}	3	°C/W

B.7.3 DC Electrical Specifications ($V_{CC} = 5.0 \text{ Vdc} \pm 5\%$)

Characteristic	Symbol	Min	Max	Unit
Input High Voltage	V_{IH}	2	V_{CC}	V
Input Low Voltage	V_{IL}	GND	0.8	V
Undershoot	—	—	0.8	V
Input Leakage Current @ 0.5–2.4 V AVEC, BCLK, BG, CDIS, IPLA, PCLK, RSTI, SCx, TBI, TLNx, TCI, TCK, TEA	I_{in}	20	20	mA
Hi-Z (Off-State) Leakage Current @ 0.5–2.4 V An, BB, CIOU, Dn, LOCK, LOCKE, R/W, SIZx, TA, TDO, TIP, TMx, TLNx, TS, TTx, UPAx	I_{TSI}	20	20	mA
Signal Low Input Current, $V_{IL} = 0.8 \text{ V}$ TMS, TDI, TRST	I_{IL}	-1.1	-0.18	mA
Signal High Input Current, $V_{IH} = 2.0 \text{ V}$ TMS, TDI, TRST	I_{IH}	-0.94	-0.16	mA
Output High Voltage, $I_{OH} = 5 \text{ mA}$	V_{OH}	2.4	—	V
Output Low Voltage, $I_{OL} = 5 \text{ mA}$	V_{OL}	—	0.5	V
Capacitance*, $V_{in} = 0 \text{ V}$, $f = 1 \text{ MHz}$	C_{in}	—	25	pF

*Capacitance is periodically sampled rather than 100% tested.

B.7.4 Power Dissipation

Frequency	Watts
Maximum Values ($V_{CC} = 5.25 \text{ V}$, $T_A = 0^\circ\text{C}$)	
20 MHz	3.2
25 MHz	3.9
33 MHz	4.9
40 MHz	TBD
Typical Values ($V_{CC} = 5 \text{ V}$, $T_A = 25^\circ\text{C}$)*	
20 MHz	2.0
25 MHz	2.4
33 MHz	3.0
40 MHz	TBD

*This information is for system reliability purposes.

B.7.5 Clock AC Timing Specifications (see Figure B-7)

Num	Characteristic	20 MHz		25 MHz		33 MHz		40 MHz		Unit
		Min	Max	Min	Max	Min	Max	Min	Max	
	Frequency of Operation	16.67	20	16.67	25	16.67	33.3	20	40	MHz
1	PCLK Cycle Time	25	30	20	30	15	30	12.5	25	nS
2	PCLK Rise Time	—	1.7	—	1.7	—	1.7	—	1.5	nS
3	PCLK Fall Time	—	1.6	—	1.6	—	1.6	—	1.5	nS
4	PCLK Duty Cycle Measured at 1.5 V	48	52	47.5	52.5	46.67	53.33	46.00	54.00	%
4a*	PCLK Pulse Width High Measured at 1.5 V	12	13	9.5	10.5	7	8	5.75	6.75	nS
4b*	PCLK Pulse Width Low Measured at 1.5 V	12	13	9.5	10.5	7	8	5.75	6.75	nS
5	BCLK Cycle Time	50	60	40	60	30	60	25	50	nS
6,7	BCLK Rise and Fall Time	—	4	—	4	—	3	—	3	nS
8	BCLK Duty Cycle Measured at 1.5 V	40	60	40	60	40	60	40	60	%
8a*	BCLK Pulse Width High Measured at 1.5 V	20	30	16	24	12	18	10	15	nS
8b*	BCLK Pulse Width Low Measured at 1.5 V	20	30	16	24	12	18	10	15	nS
9	PCLK, BCLK Frequency Stability	—	1000	—	1000	—	1000	—	1000	ppm
10	PCLK to BCLK Skew	—	9	—	9	—	N/A	—	N/A	nS

*Specification value at maximum frequency of operation.

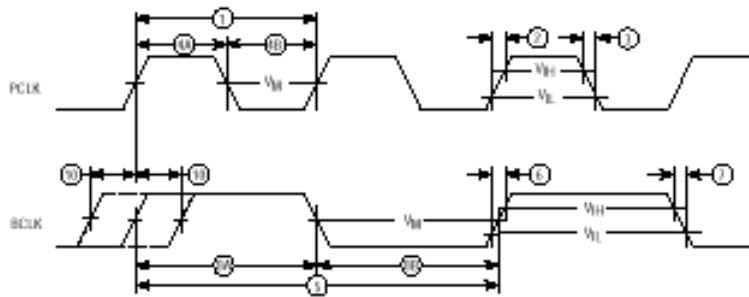


Figure B-7. Clock Input Timing Diagram

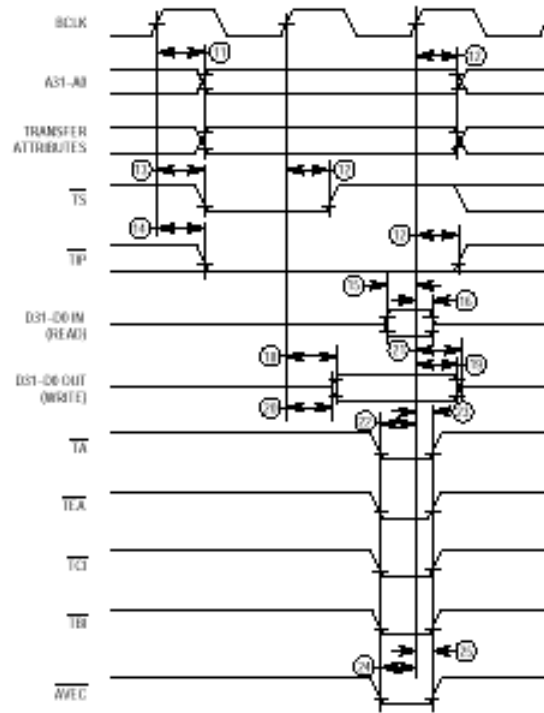
B.7.6 Output AC Timing Specifications (see Figures B-8* to B-12)

Num	Characteristic	20 MHz		25 MHz		33 MHz		40 MHz		Unit
		Min	Max	Min	Max	Min	Max	Min	Max	
11	BCLK to Address $\overline{\text{C}}\text{IOUT}$, $\overline{\text{LOCK}}$, $\overline{\text{LOCKE}}$, $\text{R}/\overline{\text{W}}$, SIZx , TLNx , TMx , TTx , UPAx Valid	11.5	35	9	30	6.5	25	5.25	24	nS
12	BCLK to Output Invalid (Output Hold)	11.5	—	9	—	6.5	—	5.25	—	nS
13	BCLK to $\overline{\text{TS}}$ Valid	11.5	35	9	30	6.5	25	5.25	24	nS
14	BCLK to $\overline{\text{TIP}}$ Valid	11.5	35	9	30	6.5	25	5.25	24	nS
18	BCLK to Data-Out Valid	11.5	37	9	32	6.5	27	5.25	26	nS
19	BCLK to Data-Out Invalid (Output Hold)	11.5	—	9	—	6.5	—	5.25	—	nS
20	BCLK to Output Low Impedance	11.5	—	9	—	6.5	—	5.25	—	nS
21	BCLK to Data-Out High Impedance	11.5	25	9	20	6.5	17	5.25	16	nS
38	BCLK to Address, $\overline{\text{C}}\text{IOUT}$, $\overline{\text{LOCK}}$, $\overline{\text{LOCKE}}$, $\text{R}/\overline{\text{W}}$, SIZx , $\overline{\text{TS}}$, TLNx , TMx , TTx , UPAx High Impedance	11.5	23	9	18	6.5	15	13	32	nS
39	BCLK to $\overline{\text{BB}}$, $\overline{\text{TA}}$, $\overline{\text{TIP}}$ High Impedance	23	33	19	28	14	25	13	—	nS
40	BCLK to $\overline{\text{BR}}$, $\overline{\text{BB}}$ Valid	11.5	35	9	30	6.5	23	5.25	14	nS
43	BCLK to $\overline{\text{MI}}$ Valid	11.5	35	9	30	6.5	25	13	19	nS
48	BCLK to $\overline{\text{TA}}$ Valid	11.5	35	9	30	6.5	25	13	34	nS
50	BCLK to $\overline{\text{IPEND}}$, $\overline{\text{PSTx}}$, $\overline{\text{RSTO}}$ Valid	11.5	35	9	30	6.5	25	5.25	14	nS

*Output timing is specified for a valid signal measured at the pin. Timing is specified driving an unterminated 30- Ω transmission line with a length characterized by a 2.5-nS one-way propagation delay. Buffer output impedance is typically 30 Ω ; the buffer specifications include approximately 5 ns for the signal to propagate the length of the transmission line and back.

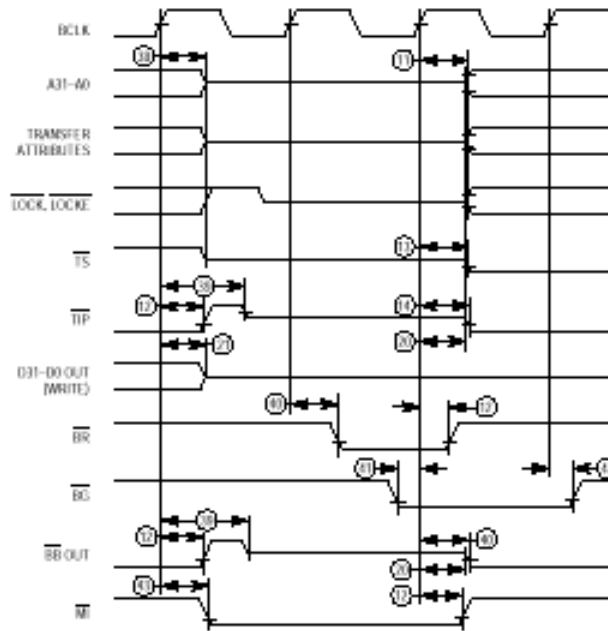
B.7.7 Input AC Timing Specifications (see Figures B-8 to B-12)

Num	Characteristic	20 MHz		25 MHz		33 MHz		40 MHz		Unit
		Min	Max	Min	Max	Min	Max	Min	Max	
15	Data-In Valid to BCLK (Setup)	6	—	5	—	4	—	3	—	nS
16	BCLK to Data-In Invalid (Hold)	5	—	4	—	4	—	3	—	nS
17	BCLK to Data-In High Impedance (Read Followed by Write)	—	61	—	49	—	36.5	—	30.25	nS
22a	\overline{TA} Valid to BCLK (Setup)	12.5	—	10	—	10	—	8	—	nS
22b	\overline{TEA} Valid to BCLK (Setup)	12.5	—	10	—	10	—	9	—	nS
22c	\overline{TCI} Valid to BCLK (Setup)	12.5	—	10	—	10	—	9	—	nS
22d	\overline{TBI} Valid to BCLK (Setup)	14	—	11	—	10	—	9	—	nS
23	BCLK to \overline{TA} , \overline{TEA} , \overline{TCI} , \overline{TBI} Invalid (Hold)	2.5	—	2	—	2	—	2	—	nS
24	\overline{AVEC} Valid to BCLK (Setup)	6	—	5	—	5	—	5	—	nS
25	BCLK to \overline{AVEC} Invalid (Hold)	2.5	—	2	—	2	—	2	—	nS
41a	\overline{BB} Valid to BCLK (Setup)	8	—	7	—	7	—	8	—	nS
41b	\overline{BG} Valid to BCLK (Setup)	10	—	8	—	7	—	2	—	nS
41c	\overline{CDIS} Valid to BCLK (Setup)	12.5	—	10	—	8	—	8	—	nS
41d	\overline{IPLA} Valid to BCLK (Setup)	5	—	4	—	3	—	3	—	nS
42	BCLK to \overline{BB} , \overline{BG} , \overline{CDIS} , \overline{IPLA} Invalid (Hold)	2.5	—	2	—	2	—	12	—	nS
44a	Address Valid to BCLK (Setup)	10	—	8	—	7	—	5	—	nS
44b	SIZx Valid to BCLK (Setup)	15	—	12	—	8	—	4	—	nS
44c	TTx Valid to BCLK (Setup)	7.5	—	6	—	8.5	—	7	—	nS
44d	R/W Valid to BCLK (Setup)	7.7	—	6	—	5	—	7	—	nS
44e	SCx Valid to BCLK (Setup)	12.5	—	10	—	11	—	8	—	nS
45	BCLK to Address SIZx, TTx, R/W, SCx Invalid (Hold)	2.5	—	2	—	2	—	3	—	nS
46	\overline{TS} Valid to BCLK (Setup)	6	—	5	—	9	—	2	—	nS
47	BCLK to \overline{TS} Invalid (Hold)	2.5	—	2	—	2	—	7	—	nS
49	BCLK to \overline{BB} High Impedance (MC68EC040 Assumes Bus Mastership)	—	11	—	9	—	9	8	—	nS
51	\overline{RSTI} Valid to BCLK	6	—	5	—	4	—	8.5	—	nS
52	BCLK to \overline{RSTI} Invalid	2.5	—	2	—	2	—	5	—	nS



NOTE: Transfer Attribute Signals = UPAs, SIZx, TTx, TMx, TLNx, RW, LOCK, LOCKE, CIOU

Figure B-8. Read/Write Timing



NOTE: Transfer Attribute Signals = UPAx, SIZx, TTx, TMx, TLNx, RW, DOUT

Figure B-9. Bus Arbitration Timing

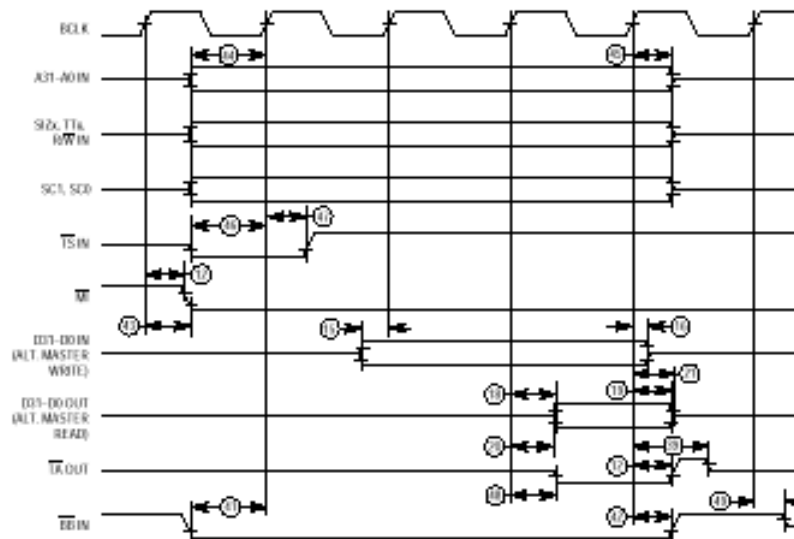


Figure B-10. Snoop Hit Timing

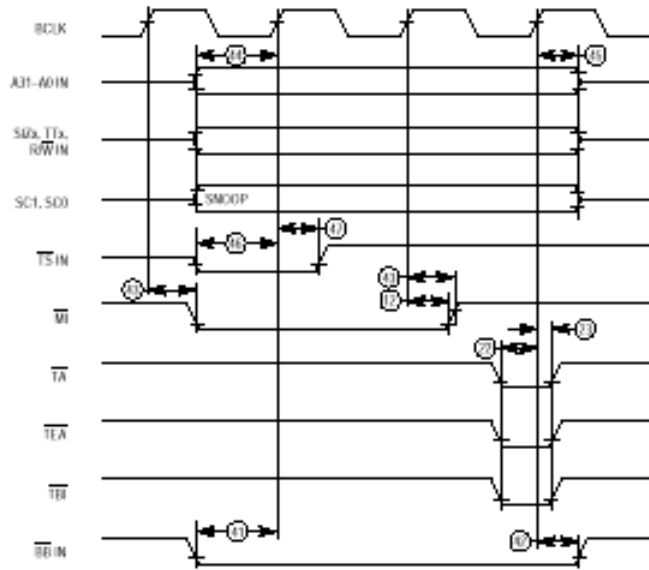


Figure B-11. Snoop Miss Timing

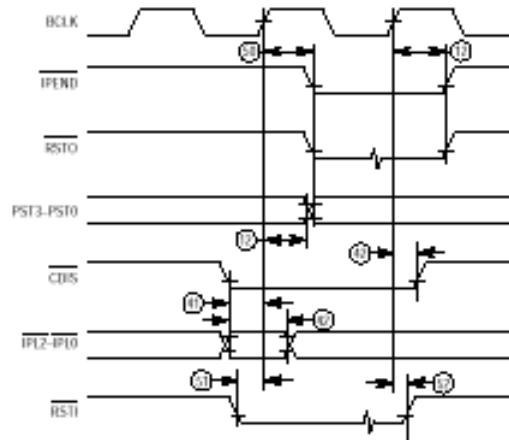


Figure B-12. Other Signal Timing